

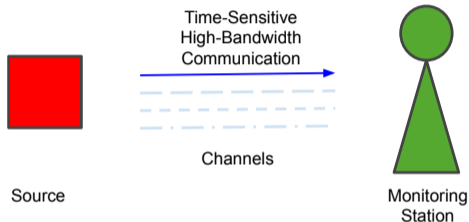
Correlated Age-of-Information Bandits

Ishank Juneja - 16D070012

Advisor: Prof. Sharayu Moharir
Department of Electrical Engineering

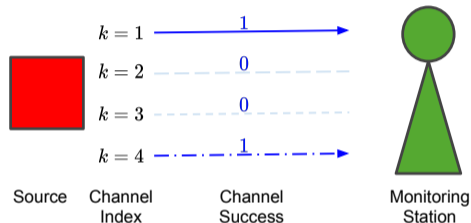
November 23, 2020

The Problem Setting



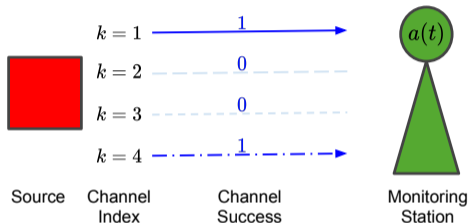
- System: A sensor node (the source) and a monitoring station
- Aim: Communicate **time-sensitive** and **high-bandwidth** information between the sensor and central monitoring station
- Problem: Find a channel selection **policy** such that cumulative performance is maximized

The Channel Selection Problem



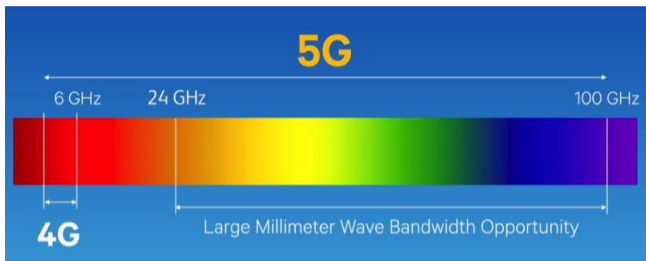
- The available bandwidth is partitioned into K frequency channels
- Schedule channel k_t among K available channels at every time step t
- Channel either successful - 1 or unsuccessful - 0
- Assume stationary channel statistics across T trials

Scheduling Problem Formulation



- Schedule in a manner that minimizes cumulative Age-of-Information
- Aol - $a(t)$ is the time elapsed since the most recent successful update
- Multi-Armed Bandit (MAB) framework applicable
- MAB policies applied and analysed in the work of [2]

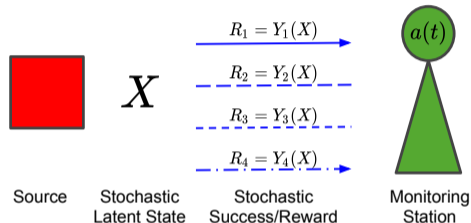
High Bandwidth - 5G Uses Shorter Waves



Link: [Image Source](#)

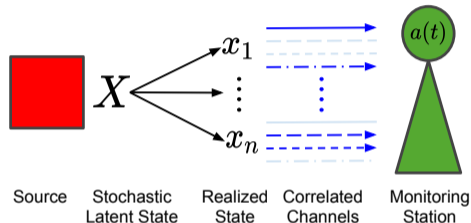
- Next generation: Higher data rates, move to higher frequency band
- 5G will use the 30GHz - 300GHz (mmWave/EHF) band
- New challenges arise: line-of-sight paths, attenuation
- Line-of-sight affects all channels, attenuation is frequency selective [3]

Performance of Channels is Correlated



- Stochastic success (reward) of arm k is $Y_k(X)$
- $Y_k(X)$ is a known deterministic function of state X
- X is a latent stochastic state with unknown distribution
- Correlation model introduced in the work of [4]

The Correlation Model



- Realizations of X lie in alphabet $\{x_1, x_2, \dots, x_n\}$
- Realization dictates which channels would be successful if used
- Successes across channels at a given time are correlated depending on the functions Y_1, Y_2, \dots, Y_K

- Variants of the UCB and Thompson Sampling policies that account for correlation analysed for the Aol regret metric
- Lower bound on Aol regret of $\Omega(\log T)$ for certain problem instances
- An upper bound on Aol regret for Correlated-UCB (CUCB) and Correlated-Thompson Sampling
- Simulations to compare the performance of UCB, Thompson Sampling and their correlated variants

Correlated Bandit Model Definitions

- Construct a Correlated Bandit instance with K arms
- Sample arm k - Obtain reward $Y_k(X)$, mean reward $\mu_k = \mathbb{E}_X[Y_k(X)]$
- Sub-optimality gap: $\Delta_k = \mu^* - \mu_k$

Definition (Expected pseudo reward and pseudo gap)

Pseudo reward for arm ℓ with respect to arm k is given by,

$$s_{\ell,k}(r) = \sup_{x: Y_k(x)=r} Y_{\ell}(x).$$

Expected pseudo reward in turn is defined as,

$$\phi_{\ell,k} = \mathbb{E}_X[s_{\ell,k}(Y_k(X))].$$

The pseudo gap is defined as $\tilde{\Delta}_{\ell,k^*} = \mu^* - \phi_{\ell,k^*}$.

Correlated Bandit Model Definitions

- If $\tilde{\Delta}_{\ell,k^*} > 0$, then arm ℓ is non-competitive
- Arm ℓ is *competitive* if $\tilde{\Delta}_{\ell,k^*} \leq 0$ and *strictly competitive* if the inequality is strict
- C denotes the number of competitive arms excluding arm k^* .

Definition (Expected pseudo reward and pseudo gap)

Pseudo reward for arm ℓ with respect to arm k is given by,

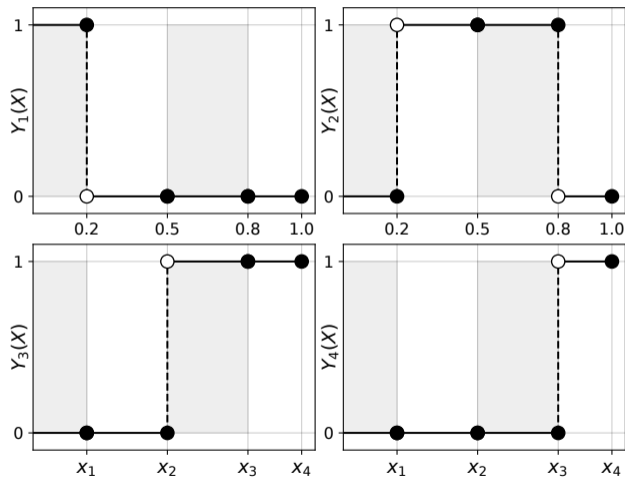
$$s_{\ell,k}(r) = \sup_{x: Y_k(x)=r} Y_{\ell}(x).$$

Expected pseudo reward in turn is defined as,

$$\phi_{\ell,k} = \mathbb{E}_X[s_{\ell,k}(Y_k(X))].$$

The pseudo gap is defined as $\tilde{\Delta}_{\ell,k^*} = \mu^* - \phi_{\ell,k^*}$.

Example 1: Correlated Bandit Model



$$\mu_1 = 1 \times 0.2 = 0.2$$

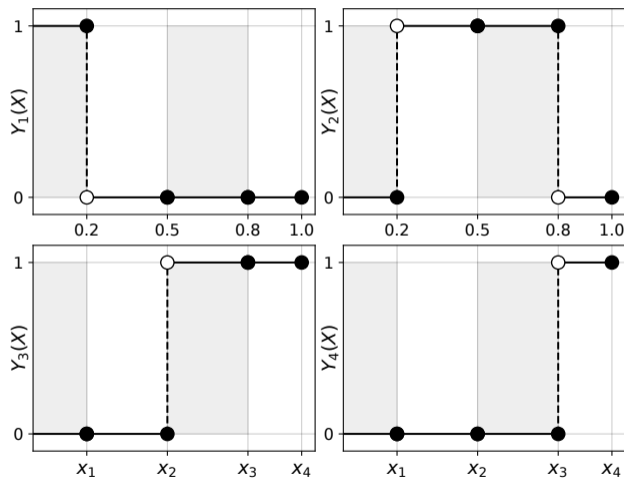
$$\mu_2 = 1 \times 0.3 + 1 \times 0.3 = 0.6$$

$$\mu_3 = 1 \times 0.3 + 1 \times 0.2 = 0.5$$

$$\mu_4 = 1 \times 0.2 = 0.2$$

- Arm 2 is optimal
- Optimal arm $k^* = 2$.
- $\mu^* = \mu_2 = 0.6$

Example 1: Correlated Bandit Model



Using Definition 1,

$$\phi_{1,2} = 1 \times 0.4 = 0.4$$

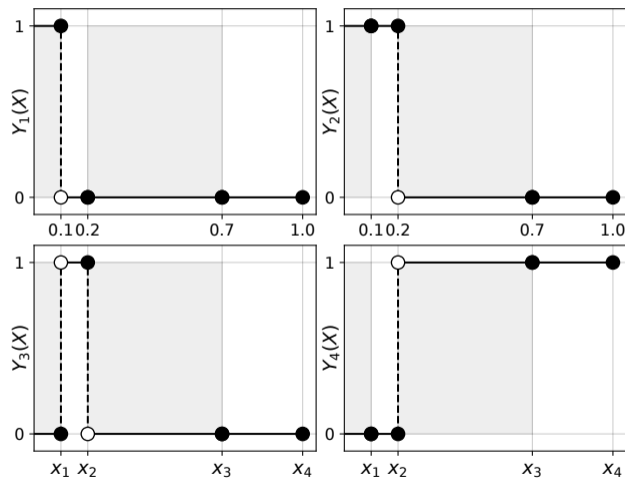
$$\phi_{2,2} = \mu_2 = 0.6$$

$$\phi_{3,2} = 1 \times 0.4 + 1 \times 0.6 = 1.0$$

$$\phi_{4,2} = 1 \times 0.4 = 0.4$$

- Only competitive sub-optimal is arm 3
- For this example $C = 1$.

Example 2: Correlated Bandit Model



$$\mu_1 = 1 \times 0.1 = 0.1$$

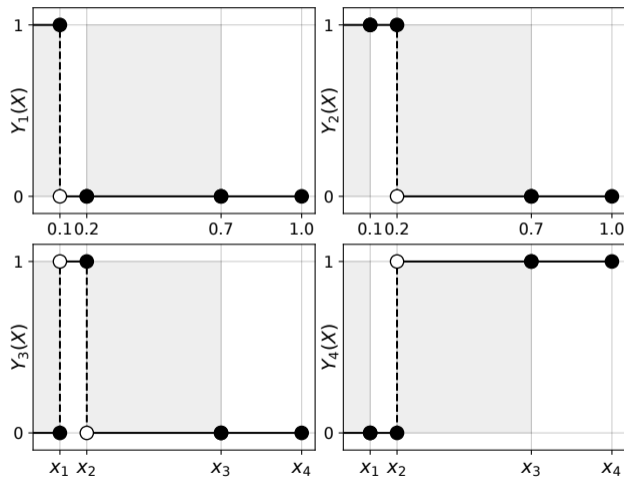
$$\mu_2 = 1 \times 0.1 + 1 \times 0.1 = 0.2$$

$$\mu_3 = 1 \times 0.1 = 0.1$$

$$\mu_4 = 1 \times 0.5 + 1 \times 0.3 = 0.8$$

- Arm 4 is optimal
- Optimal arm $k^* = 4$.
- $\mu^* = \mu_4 = 0.8$

Example 2: Correlated Bandit Model



Using Definition 1,

$$\phi_{1,4} = 1 \times 0.2 = 0.2$$

$$\phi_{2,4} = 1 \times 0.2 = 0.2$$

$$\phi_{3,4} = 1 \times 0.2 = 0.2$$

$$\phi_{4,4} = \mu_4 = 0.8$$

- No competitive sub-optimal arms
- For this example $C = 0$.

Age-of-Information (Aol) Definition

The current Aol is the time elapsed since the last successful update.

More formally,

Definition (Age-of-Information (Aol))

At the start of time slot t , let $a(t)$ denote the Aol at the central monitoring station and let $u(t)$ denote the time index at which the recent most successful update was received by the monitoring station. Then, $a(t) = t - u(t)$. Alternatively,

$$a(t) = \begin{cases} 1 & \text{if the update in } t - 1 \text{ succeeds} \\ a(t - 1) + 1 & \text{otherwise.} \end{cases}$$

Definition (Age-of-Information Regret (Aol Regret))

Aol regret for a policy ρ , over T slots is given by,

$$R_\rho(T) = \sum_{t=1}^T \mathbb{E}[a_\rho(t) - a^*(t)] = \sum_{t=1}^T \mathbb{E}[a_\rho(t)] - \frac{T}{\mu^*}, \quad (1)$$

where (1) follows from the expectation of a geometric random variable with parameter μ^* .

- $a_\rho(t)$ denotes the Aol in time slot t under policy ρ .
- $a^*(t)$ is the Aol under the optimal policy.

Aol Regret Lower Bound - Policy Family

Lower Bound on Aol regret is derived for a certain α -consistent family of policies.

Definition (α -consistent policies [6])

Let k_s denote the index of the channel scheduled in time-slot s . The index k^* denotes the index of the optimal channel. A scheduling policy is called α -consistent, for a constant $\alpha \in (0, 1)$, if there exists an instance dependent constant M such that,

$$\mathbb{E} \left[\sum_{s=1}^t \mathbb{1}\{k_s = k\} \right] \leq Mt^\alpha, \quad \forall k \neq k^*. \quad (2)$$

Aol Regret Lower Bound

Theorem (Lower bound on Aol regret)

If a bandit instance I has at least one strictly competitive arm k with $\tilde{\Delta}_{k,k^*} < 0$, then for any α -consistent policy ρ , we have,

$$R_\rho(T) \geq \max_{k \in \mathcal{C}'} \frac{\Delta_k}{D(P_k, P'_k)} \frac{(1 - \alpha) \log T - \log(4M)}{\mu^*}.$$

Otherwise, if $\tilde{\Delta}_{k,k^*} \geq 0 \forall k \in [K]$, $R_\rho(T) \geq 0$.

- $D(P_k, P'_k)$ is the KL divergence between the reward distribution of arm k and a suitably chosen perturbed reward distribution
- \mathcal{C}' is the set of strictly competitive arms
- M is an instance dependent constant as in Definition 4

Definitions Associated With CUCB

Let C denote the number of competitive sub-optimal arms and \mathcal{C} denote the set of competitive arms inclusive of k^* .

$$t_o = \inf \left\{ \tau \geq 2 : \Delta_{\min}, \tilde{\Delta}_{k,k^*} \geq 4\sqrt{\frac{2K \log \tau}{\tau}} \right\},$$

$$U_{k,\text{CUCB}}^{(nc)} = Kt_0 + K^3 \sum_{t=Kt_0}^T 2\left(\frac{t}{K}\right)^{-2} + \sum_{t=1}^T 3t^{-3},$$

$$U_{k,\text{CUCB}}^{(c)} = 8\frac{\log(T)}{\Delta_k^2} + \left(1 + \frac{\pi^2}{3}\right) + \sum_{t=1}^T 2Kt \exp\left(-\frac{t\Delta_{\min}^2}{2K}\right).$$

Where, $\Delta_{\min} = \min_{k \neq k^*} \mu^* - \mu_k$.

Upper bound on Aol Regret Under CUCB

Theorem (Upper bound on Aol regret under CUCB)

Let $\mu_{\min} = \min_k \mu_k$, then for $T > t_0$,

$$\begin{aligned}\mathbb{E}[R_{\text{CUCB}}(T)] &\leq \frac{1 - \mu^*}{\mu^* \mu_{\min}} + \left(\frac{1}{\mu_{\min}} - \frac{1}{\mu^*} \right) \left(\sum_{k' \in [K] \setminus \mathcal{C}} \Delta_{k'} U_{k, \text{CUCB}}^{(nc)} + \sum_{k \in \mathcal{C} \setminus \{k^*\}} \Delta_k U_{k, \text{CUCB}}^{(c)} \right) \\ &= O(1) + O(C \log T),\end{aligned}$$

and for $T \leq t_0$,

$$\mathbb{E}[R_{\text{CUCB}}(T)] \leq \left(\frac{1}{\mu_{\min}} - \frac{1}{\mu^*} \right) T.$$

Definitions Associated With C - Thompson Sampling

Let C denote the number of competitive sub-optimal arms and \mathcal{C} denote the set of competitive arms inclusive of k^* .

$$t_b = \inf \left\{ \tau \geq \exp(11\beta) : \Delta_{\min}, \tilde{\Delta}_{k,k^*} \geq 6\sqrt{\frac{2K\beta \log \tau}{\tau}} \right\},$$

$$U_{k,\text{CTS}}^{(nc)} = Kt_b + \sum_{t=1}^T 3t^{-3} + K^2 \sum_{t=Kt_b}^T \left((2K+3) \left(\frac{t}{K}\right)^{-2} + \left(\frac{t}{K}\right)^{1-2\beta} \right),$$

$$U_{k,\text{CTS}}^{(c)} = 18 \frac{\log(T\Delta_k^2)}{\Delta_k^2} + \exp(11\beta) + \frac{9}{\Delta_k^2} + \sum_{t=1}^T 2Kt \exp\left(-\frac{t\Delta_{\min}^2}{2K}\right).$$

Where, $\Delta_{\min} = \min_{k \neq k^*} \mu^* - \mu_k$ and $\beta > 1$ is a parameter of the Thompson Sampling with Gaussian priors algorithm.

Upper Bound on Aol Regret Under C - Thompson Sampling

Theorem (Upper bound on Aol regret under CTS)

Then, for any choice of $\beta > 1$ and for $T > t_b$,

$$\begin{aligned}\mathbb{E}[R_{\text{CTS}}(T)] &\leq \frac{1 - \mu^*}{\mu^* \mu_{\min}} + \left(\frac{1}{\mu_{\min}} - \frac{1}{\mu^*} \right) \left(\sum_{k' \in [K] \setminus \mathcal{C}} \Delta_{k'} U_{k, \text{CTS}}^{(nc)} + \sum_{k \in \mathcal{C} \setminus \{k^*\}} \Delta_k U_{k, \text{CTS}}^{(c)} \right) \\ &= O(1) + O(C \log T),\end{aligned}$$

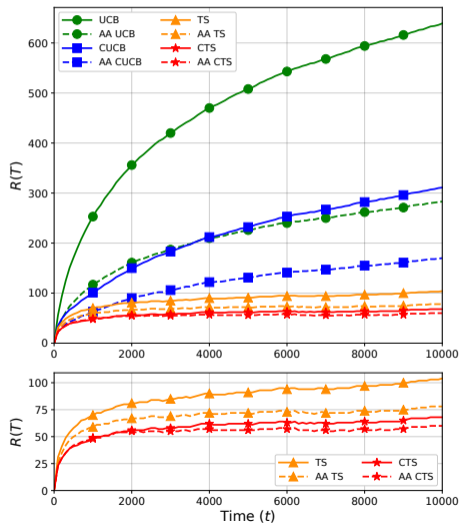
and for $T \leq t_b$,

$$\mathbb{E}[R_{\text{CTS}}(T)] \leq \left(\frac{1}{\mu_{\min}} - \frac{1}{\mu^*} \right) T.$$

Discussion About Regret Bounds

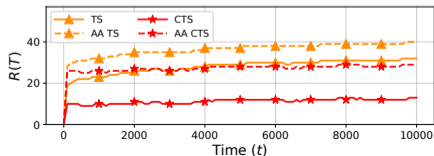
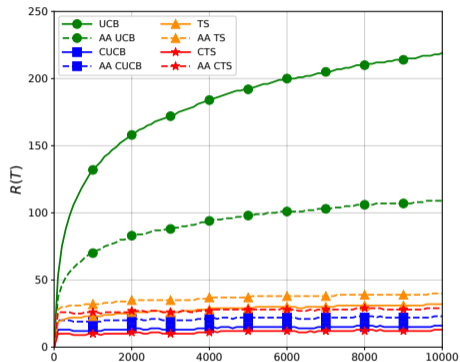
- If $C > 0$ for a Correlated Bandit instance, and if at least one arm is strictly competitive, then the lower bound and upper bound on Aol regret are both $O(\log T)$
- When there are no competitive arms ($C = 0$), there is no meaningful lower bound on the expected Aol regret
- The $C = 0$ case agrees with the fact that the set $\mathcal{C} \setminus \{k^*\}$ being empty results in an $O(1)$ upper bound on Aol regret
- For these cases of Correlated Bandit instances Aol regret bounds are order-optimal

Simulation Results for Example 1



- An Aol-aware policy follows the policy or is greedy based on a Threshold [2]
- CTS and its Aol-aware variant perform the best on Aol regret
- CUCB and CTS have significantly lower Aol regret compared to UCB and TS

Simulation Results for Example 2



- Bandit instance in Example 2 had no competitive sub-optimal arms, i.e. $C = 0$
- As predicted by the bounds in the preceding Theorems both CUCB and CTS have constant expected Aol regret
- The Aol-aware variant of a policy need not perform better than its parent policy as is the case for CUCB, CTS and TS in this example

Key Takeaways and Conclusion

- 5G frequency band means the problem of line-of-sight occlusion becomes more significant
- Exploit correlation to identify certain channels as sub-optimal within a few time steps
- Assumption: Deterministic reward functions
- Strength: Once determined, reward functions can be applied in other communication systems with a similar configuration but a different and unknown distribution of X
- Distribution agnostic model and algorithms analysed in this work would be highly beneficial in such scenarios

Directions Planned to be Pursued

- Dropping the assumption of channel reward distribution being stationary across time
- Theoretical analysis of Aol-aware policies
- Alternate correlation models to exploit 0-1 Binary rewards
- Aol regret upper bound for C - Thompson Sampling with Beta priors

References I

- [1] I. Juneja, S. Fatale, and S. Moharir, “Correlated age-of-information bandits,” *arXiv preprint arXiv:2011.05032*, 2020.
- [2] S. Fatale, K. Bhandari, U. Narula, S. Moharir, and M. K. Hanawal, “Regret of age-of-information bandits,” *arXiv preprint arXiv:2001.09317*, 2020.
- [3] K. C. Huang and Z. Wang, *Millimeter Wave Characteristics*. John Wiley & Sons, Ltd, 2011, ch. 1, pp. 1–31.
- [4] S. Gupta, S. Chaudhari, G. Joshi, and O. Yağan, “Multi-armed bandits with correlated arms,” *arXiv preprint arXiv:1911.03959*, 2019.
- [5] S. Kaul, R. Yates, and M. Gruteser, “Real-time status: How often should one update?” in *2012 Proceedings IEEE INFOCOM*, March 2012, pp. 2731–2735.
- [6] T. L. Lai and H. Robbins, “Asymptotically efficient adaptive allocation rules,” *Advances in applied mathematics*, vol. 6, no. 1, pp. 4–22, 1985.
- [7] W. R. Thompson, “On the likelihood that one unknown probability exceeds another in view of the evidence of two samples,” *Biometrika*, vol. 25, no. 3/4, pp. 285–294, 1933.

References II

- [8] L. Zhou, “A survey on contextual multi-armed bandits,” 2015.
- [9] V. Dani, T. Hayes, and S. Kakade, “Stochastic linear optimization under bandit feedback.” 01 2008, pp. 355–366.
- [10] A. Kosta, N. Pappas, V. Angelakis *et al.*, “Age of information: A new concept, metric, and tool,” *Foundations and Trends® in Networking*, vol. 12, no. 3, pp. 162–259, 2017.
- [11] V. Tripathi and S. Moharir, “Age of Information in Multi-Source Systems,” in *GLOBECOM 2017-2017 IEEE Global Communications Conference*. IEEE, 2017, pp. 1–6.
- [12] P. R. Jhunjunwala and S. Moharir, “Age-of-Information aware scheduling,” *SPCOM*, 2018.
- [13] I. Kadota, A. Sinha, and E. Modiano, “Optimizing age of information in wireless networks with throughput constraints,” in *Proc. INFOCOM*, 2018.
- [14] B. Sombabu and S. Moharir, “Age-of-Information Aware Scheduling for Heterogeneous Sources,” in *Proceedings of the 24th Annual International Conference on Mobile Computing and Networking*, ser. MobiCom '18. New York, NY, USA: ACM, 2018, pp. 696–698.

- [15] A. B. Tsybakov, *Introduction to Nonparametric Estimation*, ser. Springer series in statistics. Springer, 2009.
- [16] T. Lattimore and C. Szepesvári, *Bandit Algorithms*. Cambridge University Press, 2020.

End of Slides

Thank You!

Appendix: CUCB Decision Making Algorithm [4] - I

- 1: **Input:** Pseudo-rewards $s_{\ell,k}(r)$
- 2: **Initialize:** Set $\hat{\mu}_k, \hat{\phi}_{\ell,k}$ and n_k as $0 \forall k \in [K]$.
- 3: **while** $1 \leq t \leq K$ **do**
- 4: Schedule update on Channel $k_t = t$
- 5: Receive reward r_t drawn from $\text{Ber}(\mu_{k_t})$
- 6: $\hat{\mu}_{k_t} = r_t$
- 7: $n_{k_t}(t) = 1$
- 8: $t = t + 1$
- 9: **end while**
- 10: **while** $t \geq K + 1$ **do**
- 11: Find $\mathcal{S}_t = \{k : n_k(t-1) \geq \frac{t-1}{K}\}$, the set of arms pulled a significant number of times till $t-1$. Define $k^{\text{emp}}(t) = \arg \max_{k \in \mathcal{S}_t} \hat{\mu}_{k_t}$
- 12: Initialize the empirically competitive set \mathcal{A}_t as $\{\}$
- 13: **for** $k \in [K]$ **do**

Appendix: CUCB Decision Making Algorithm [4] - II

```
14:     if  $\min_{\ell \in \mathcal{S}_t} \hat{\phi}_{k,\ell}(t) \geq \hat{\mu}_{k^{\text{emp}}}(t)$  then
15:         Add empirically competitive arms  $k$  to the set:  $\mathcal{A}_t = \mathcal{A}_t \cup \{k\}$ 
16:     end if
17: end for
18: Schedule update on Channel  $k_t$  such that,  $k_t = \arg \max_{k \in \mathcal{A}_t \cup \{k^{\text{emp}}(t)\}} \hat{\mu}_{k_t} + \sqrt{\frac{2 \log t}{n_{k_t}(t-1)}}$ 
19: Receive reward  $r_t$  drawn from  $\text{Ber}(\mu_{k_t})$ 
20:  $\hat{\mu}_{k_t} = (\hat{\mu}_{k_t} \cdot n_{k_t}(t-1) + r_t) / (n_{k_t}(t-1) + 1)$ 
21:  $n_{k_t}(t) = n_{k_t}(t-1) + 1$ 
22:  $\hat{\phi}_{k,k_t} = \sum_{\tau: k_\tau = k_t} s_{k,k_\tau}(r_\tau) / n_{k_t}(t) \quad \forall k \neq k_t$ 
23:  $t = t + 1$ 
24: end while
```

Appendix: C-TS Decision Making Algorithm [4] - I

- 1: **Input:** Pseudo-rewards $s_{\ell,k}(r)$
- 2: **Initialize:** Set the number of successes, $S_k(t)$, failures, $F_k(t)$, $\hat{\mu}_k$, $\hat{\phi}_{\ell,k}$ and n_k as 0 $\forall k \in [K]$.
- 3: **while** $t \geq 1$ **do**
- 4: Find $\mathcal{S}_t = \{k : n_k(t-1) \geq \frac{t-1}{K}\}$, the set of arms pulled a significant number of times till $t-1$. Define $k^{\text{emp}}(t) = \arg \max_{k \in \mathcal{S}_t} \hat{\mu}_{k_t}$
- 5: Initialize the empirically competitive set \mathcal{A}_t as $\{\}$
- 6: **for** $k \in [K]$ **do**
- 7: **if** $\min_{\ell \in \mathcal{S}_t} \hat{\phi}_{k,\ell}(t) \geq \hat{\mu}_{k^{\text{emp}}}(t)$ **then**
- 8: Add empirically competitive arms k to the set: $\mathcal{A}_t = \mathcal{A}_t \cup \{k\}$
- 9: **end if**
- 10: **end for**
- 11: For each k in $[K]$, draw a sample $\theta_k(t)$, where,
 $\theta_k(t) \sim \text{Beta}(S_k(t-1) + 1, F_k(t-1) + 1)$

- 12: Schedule update on Channel k_t such that $k_t = \arg \max_{k \in \mathcal{A}_t \cup \{k^{\text{emp}}(t)\}} \theta_k(t)$
- 13: Receive reward r_t drawn from $\text{Ber}(\mu_{k_t})$
- 14: $S_{k_t}(t) = S_{k_t}(t-1) + r_t$
- 15: $F_{k_t}(t) = F_{k_t}(t-1) + (1 - r_t)$
- 16: $\hat{\mu}_{k_t} = (\hat{\mu}_{k_t} \cdot n_{k_t}(t-1) + r_t) / (n_{k_t}(t-1) + 1)$
- 17: $n_{k_t}(t) = n_{k_t}(t-1) + 1$
- 18: $\hat{\phi}_{k,k_t} = \sum_{\tau: k_\tau = k_t} s_{k,k_\tau}(r_\tau) / n_{k_t}(t) \forall k \neq k_t$
- 19: $t = t + 1$
- 20: **end while**